

Internet voting by CHvote

@SwissCyberStorm
18th october 2017



REPUBLIQUE
ET CANTON
DE GENEVE

POST TENEBRAS LUX

Département de la sécurité et de l'économie
Direction générale des systèmes d'information

Short Bio

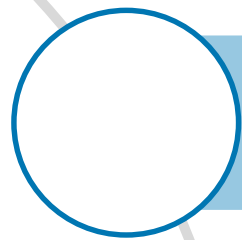
Thomas Hofer / [@thhofer](#) / thomas.hofer@etat.ge.ch

- EPFL MSc in IT
- Master thesis @ CERN Computer Security Team
- IT / Java consultant

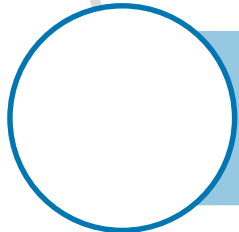
- Now
 - Java DEV & AppSec
 - Internet voting cryptography @ State of Geneva

- Outside from work
 - OWASP-Geneva co-chapter leader
 - Married, 2 kids

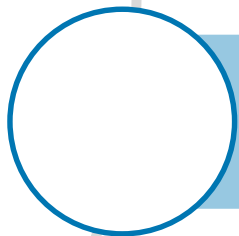
Outline



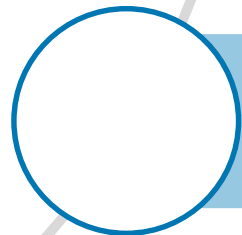
CHvote



Control components: in theory

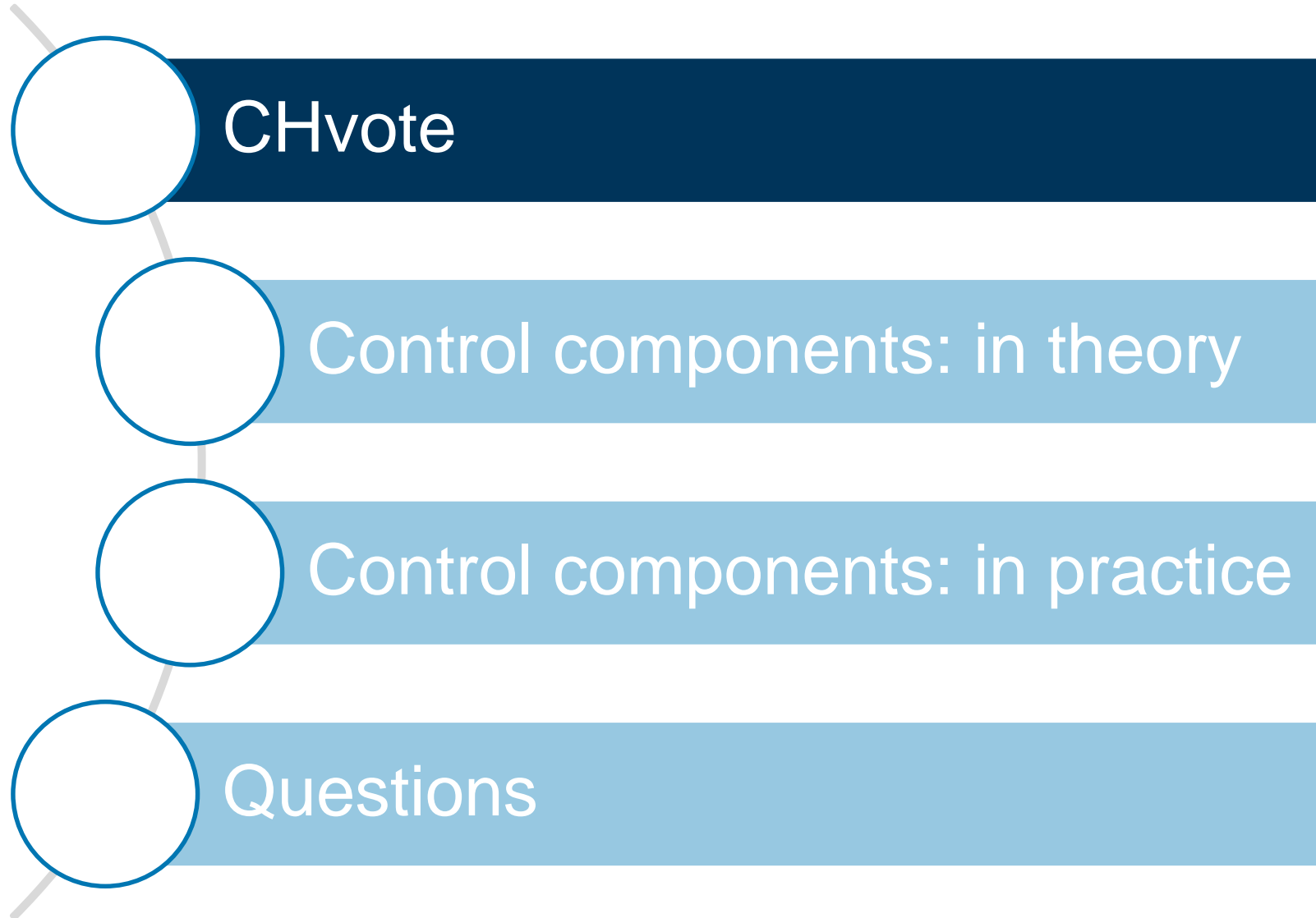


Control components: in practice



Questions

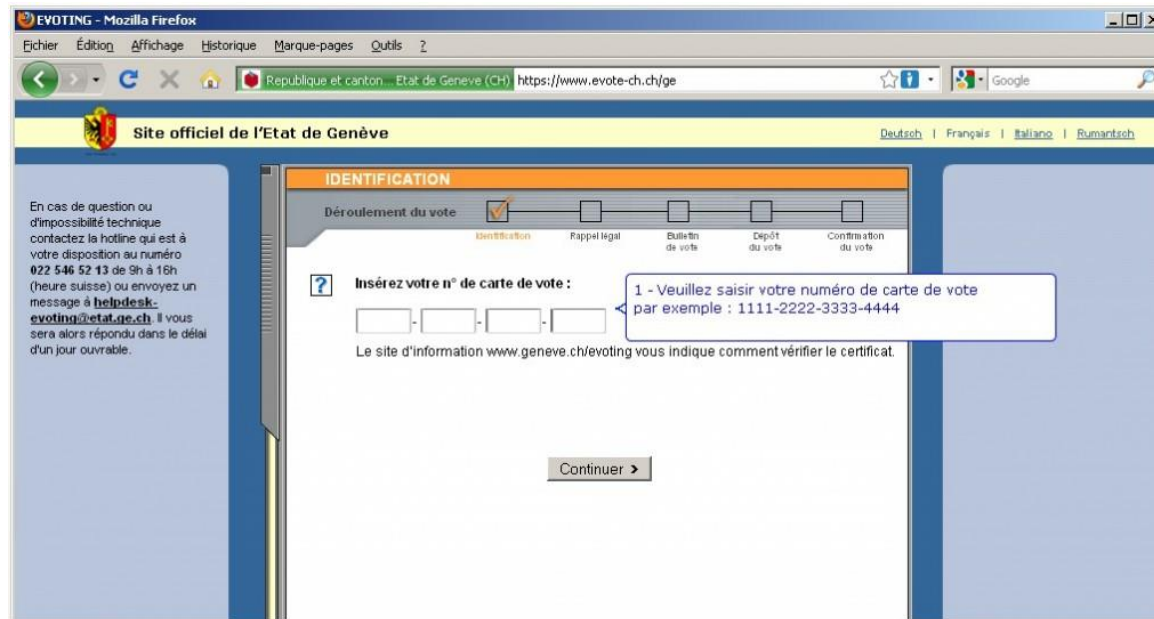
Outline



The past of CHvote

First generation E-Voting system

- 2001: start of project
- 2003: first use



- Partners



The present of CHvote

Individual verifiability & major appearance overhaul

The screenshot shows the 'Vote Electronique' interface for the Canton of Geneva. It features a progress bar with six steps: 1. Identification, 2. Rappel légal, 3. Bulletin de vote (highlighted), 4. Récapitulatif, 5. Vérification, and 6. Finalisation du vote. Below the progress bar, there is a section titled 'ETAPE 3: BULLETIN DE VOTE' with instructions to answer questions. Three federal questions are listed, each with 'OUI' and 'NON' radio buttons. To the right, there is a 'Contact' section with a phone number and an email address, and a 'FAQ du vote électronique' section with three questions.



Liste de codes pour la carte n° 5874-8863-1400-8743

Votation fédérale

Question 1

Acceptez-vous l'arrêté fédéral du 20 juin 2013 portant règlement du financement et de l'aménagement de l'infrastructure ferroviaire (Contre-projet direct à l'initiative populaire "Pour les transports publics", qui a été retirée) ?

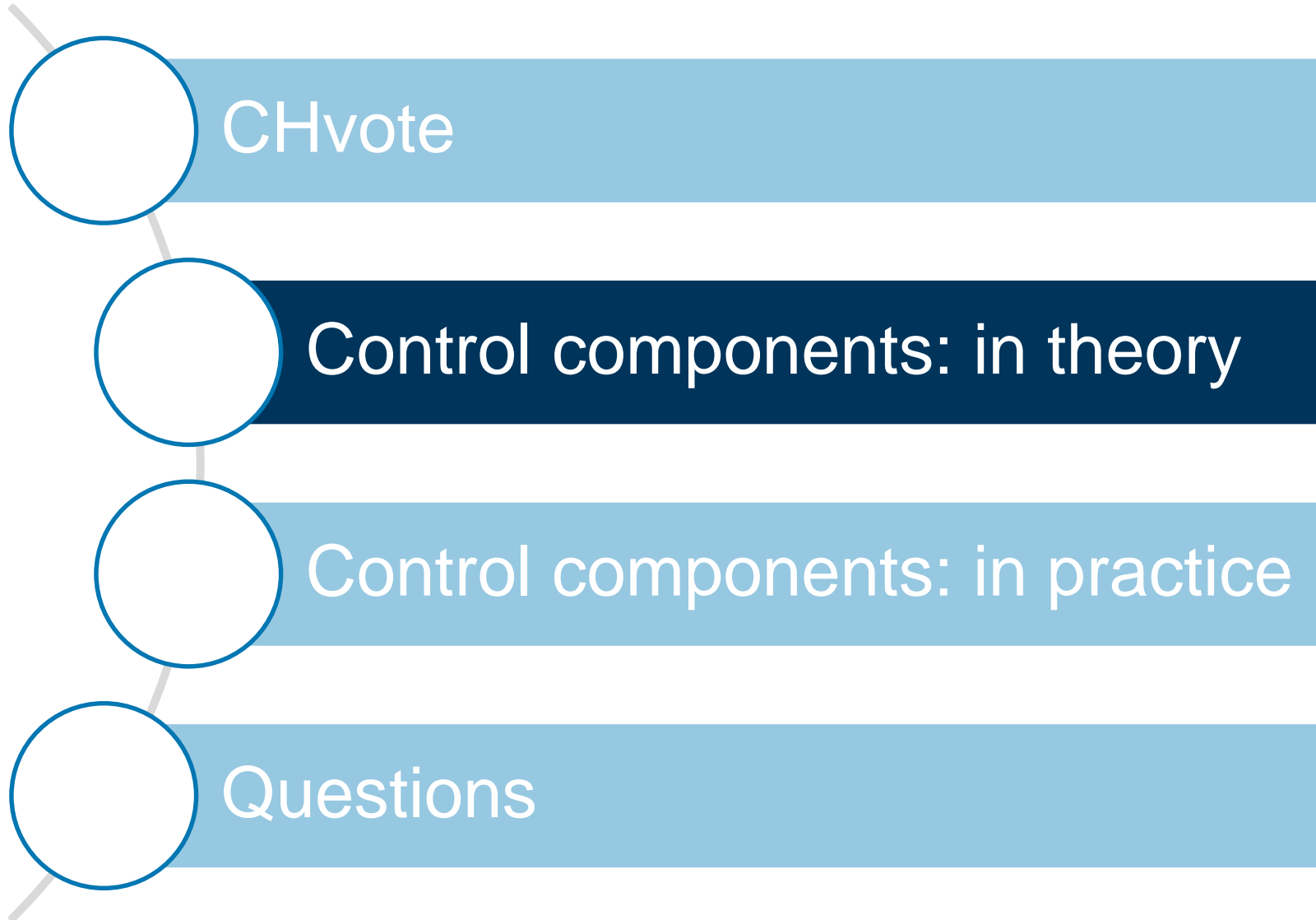
Oui	Non	Blanc
A2B4	J5B9	Z8H5

The future of CHvote

End-to-end verifiable internet voting protocol

- New academic partnerships
 - Berner Fachhochschule
 - INRIA / Bristol
 - ITU Copenhagen
 - ...
- New cryptographic protocol
 - End-to-end encryption
 - Universal Verifiability
 - Control Components
- Currently in development, ETA: 2019

Outline



Federal requirements

New Ordinance on Electronic Voting

- Published in 2013, enacted 2014
 - Collaborative work between lawmakers, academia and operating staff
- Compliance levels
 - The higher the compliance, the more voters allowed
- Reference
 - <https://www.bk.admin.ch/themen/pore/evoting/07979/index.html>

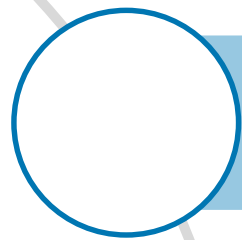
Federal requirements

Control Components

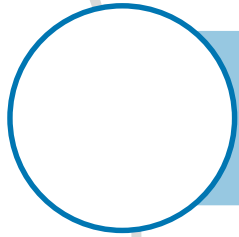
The trustworthy part of the system includes either one or a small number of groups of independent components secured by special measures (control components). Their use must also make any abuse recognisable if per group only one of the control components works correctly and in particular is not manipulated unnoticed. – *VEleS, art. 5, par. 6*

- i.e. trust splitting in "anytrust" mode: only 1-of-k needs to be honest, all k need to be available
- Control components must use different OS

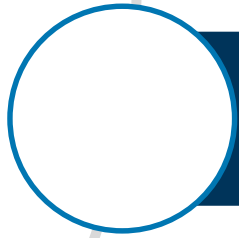
Outline



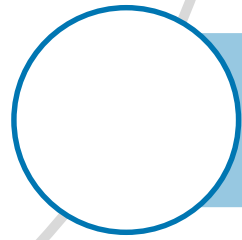
CHvote



Control components: in theory



Control components: in practice



Questions

Security role of control components

A brief overview

- Homomorphic encryption (El Gamal)
- Public / private credentials per voter
 - Each Control Component holds a share
- 1 key pair per Control Component
 - "system" public key is the product of public keys
- 1 ballot-box shuffle per Control Component

Voter authentication

In the context of trust splitting

- Private credentials:
 - Signed and encrypted for printing authority
 - Used for signing the ballot
- Public credentials:
 - Shared and combined
 - Each Control Component must verify credentials

Ballot encryption

How ?

- Voting client uses "system" key to encrypt ballot
 - "system" key is product of each CC's public key
- All Control Components are required for decryption
- (partial) Decryption by CC is not allowed until:
 - Ballot box has been shuffled by all CCs
 - Each shuffle has been verified to be valid

Oblivious Transfer

What does it mean and why is it useful?

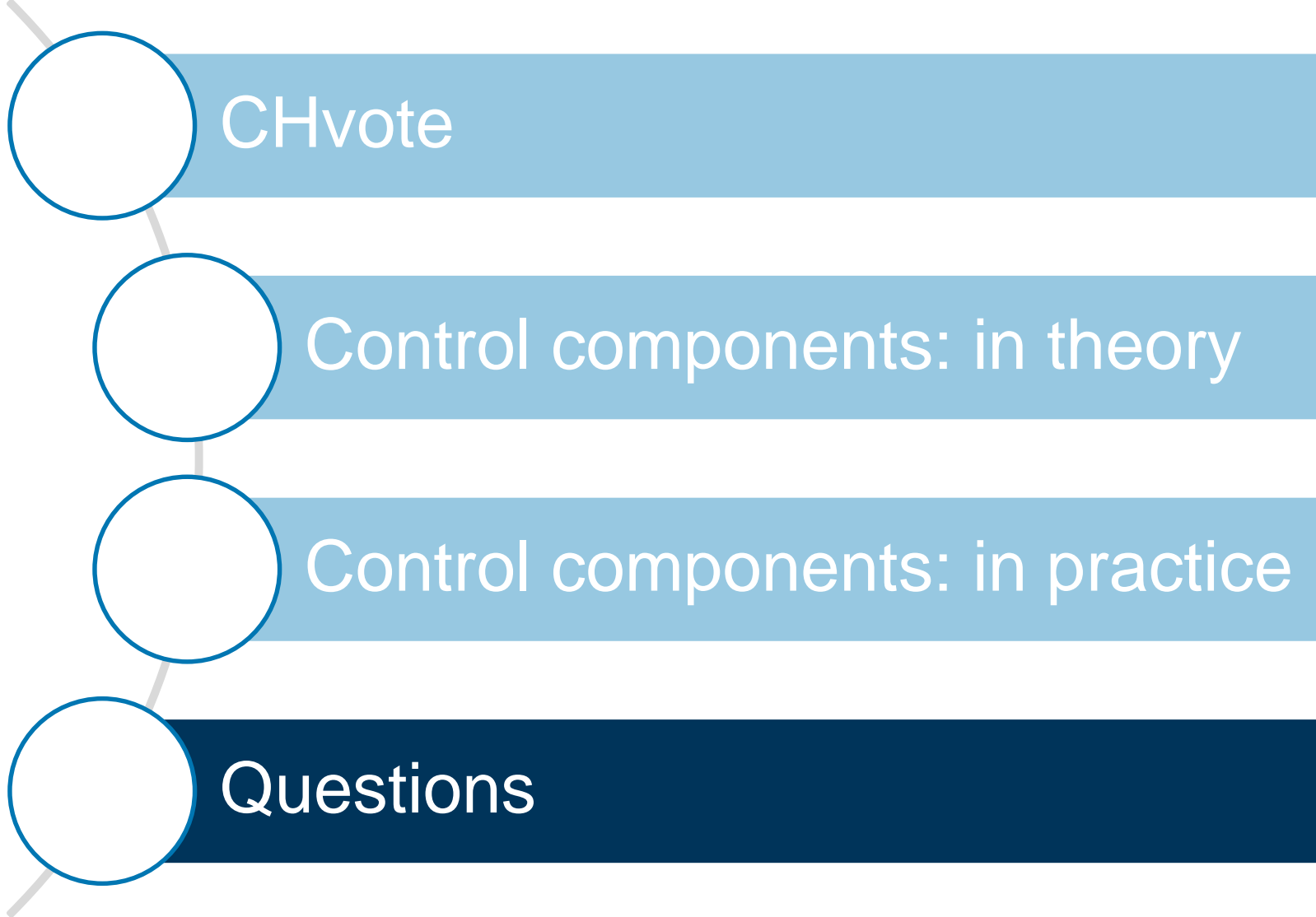
- In short
 - Server knows n secret messages
 - Client allowed to retrieve k secret messages
 - Server cannot know which messages the client asked for
 - *Enables individual verifiability combined with end-to-end encryption*
- In detail
 - [Cast-as-Intended Verification in Electronic Elections Based on Oblivious Transfer](#)

Ballot-box shuffling

Why is it complicated and how does it work

- Re-encrypting mix-net
 - Each CC **re-encrypts** each ballot and **shuffles** them
- Shuffled → simple pre-image proofs would not work
- Re-encrypted → ciphertexts are not equal
- Need for a specific proof that the cryptographic shuffle is valid

Outline





Further reading

And references

- Published protocol specification
→ <https://eprint.iacr.org/2017/325>
- Published PoC code
→ <https://github.com/republique-et-canton-de-geneve/chvote-protocol-poc>
- Federal requirements
→ <https://www.bk.admin.ch/themen/pore/evoting/07979/index.html>

Thank you!



Thomas Hofer



thomas.hofer@etat.ge.ch



@thhofer



REPUBLIQUE
ET CANTON
DE GENEVE

POST TENEBRAS LUX

Département de la sécurité et de l'économie
Direction générale des systèmes d'information